



Temporal  
cores

V. Batagelj,  
M. Cerinšek

Definitions

Algorithms

Results

Conclusions

References

# Temporal cores in networks

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**XXXVIII Sunbelt 2018**

Utrecht, The Netherlands, June 26 – July 1, 2018



# Outline

## Temporal cores

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- 1 Definitions
- 2 Algorithms
- 3 Results
- 4 Conclusions
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Current version of slides (June 30, 2018 at 04:44): [Sunbelt'18 slides PDF](#)



# Cores

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Network:  $\mathcal{N} = (\mathcal{V}, \mathcal{L}, \mathcal{P}, \mathcal{W})$ ;

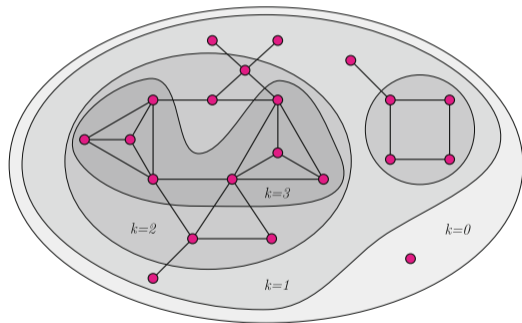
$n = |\mathcal{V}|, m = |\mathcal{L}|$

A subgraph  $\mathcal{H} = (\mathcal{C}, \mathcal{L}(\mathcal{C}))$  induced by the set of nodes  $\mathcal{C}$  is a ***k*-core** or a core of order  $k$  iff  $\forall v \in \mathcal{C} : \deg_{\mathcal{H}}(v) \geq k$  and  $\mathcal{H}$  is the maximum subgraph with this property (Seidman 1983).

The core of maximum order – main core.

The core number of node  $v$  is the highest order of a core that contains this node.

Batagelj and Zaveršnik (2003, 2011) proposed a very fast algorithm for determining core numbers.





# Generalized cores

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Node property function:  $p(v, \mathcal{C}); v \in \mathcal{V}, \mathcal{C} \subseteq \mathcal{V}, p : L \rightarrow \mathbb{R}^+$ .

Properties:

$p(v, \mathcal{C})$  is *local*:  $p(v, \mathcal{C}) = p(v, N(v, \mathcal{C})) \forall v \in \mathcal{V}$

$p(v, \mathcal{C})$  is *monotone*:  $\mathcal{C}_1 \subset \mathcal{C}_2 \Rightarrow \forall v \in \mathcal{V} : p(v, \mathcal{C}_1) \leq p(v, \mathcal{C}_2)$ .

The subgraph  $\mathcal{H} = (\mathcal{C}, \mathcal{L}(\mathcal{C}))$  induced by the set  $\mathcal{C} \subseteq \mathcal{V}$  is a *p-core* at level  $t \in \mathbb{R}$  iff  $\forall v \in \mathcal{C} : t \leq p(v, \mathcal{C})$  and  $\mathcal{C}$  is a maximal such set.



# Generalized cores

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## Examples of node property function

- 1  $p_1(v, \mathcal{C}) = \text{deg}_{\mathcal{C}}(v)$ : node degree within  $\mathcal{C}$
- 2  $p_2(v, \mathcal{C}) = \text{indeg}_{\mathcal{C}}(v) + \text{outdeg}_{\mathcal{C}}(v)$ : if lines are directed it holds  $p_2 = p_1$
- 3  $p_3(\mathbf{v}, \mathcal{C}) = \sum_{u \in N(\mathbf{v}, \mathcal{C})} \mathbf{w}(\mathbf{v}, u)$  for  $\mathbf{w} : L \rightarrow \mathbb{R}_0^+$  : **sum of weights of incident lines within  $\mathcal{C}$**
- 4  $p_4(v, \mathcal{C}) = \max_{u \in N(v, \mathcal{C})} w(v, u)$  for  $w : L \rightarrow \mathbb{R}$ : maximal weight of incident lines within  $\mathcal{C}$
- 5  $p_5(v, \mathcal{C}) = \frac{\text{deg}_{\mathcal{C}}(v)}{\text{deg}(v)}$  if  $\text{deg}(v) > 0$  else  $f_5(v, \mathcal{C}) = 0$ : fraction of neighbors within  $\mathcal{C}$ .
- 6  $p_6(v, \mathcal{C}) = \frac{\sum_{u \in N(v, \mathcal{C})} w(v, u)}{\sum_{u \in N(v)} w(v, u)}$  for  $w : L \rightarrow \mathbb{R}_0^+$ : fraction of sum of weights of incident lines within  $\mathcal{C}$ .



# Temporal network

## Temporal cores

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A *temporal network*

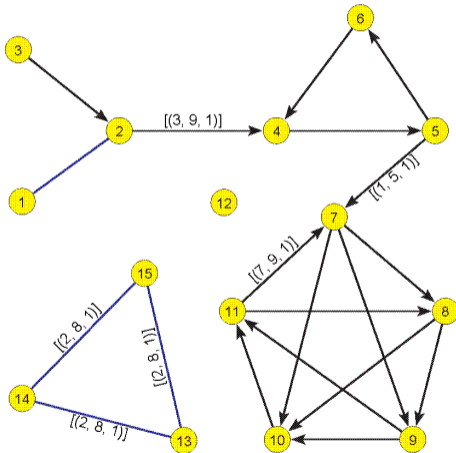
$$\mathcal{N}_{\mathcal{T}} = (\mathcal{V}, \mathcal{L}, \mathcal{T}, \mathcal{P}, \mathcal{W})$$

is obtained by attaching the time  $\mathcal{T}$  to an ordinary network, where  $\mathcal{T}$  is a set of time points:  $t \in \mathcal{T}$  which are usually integers or reals.

Temporal quantities (TQ) are assigned to nodes and links:

a TQ is a list of triples  $(s, f, v)$ :  $s$  - start,  $f$  - finish of time interval  $[s, f)$ ,  $v$  - value.

$T(v)$  - the activity set of time points for the node  $v$ ;  $T(l)$  the activity set of time points for the link  $l$





# Sum and product of temporal quantities

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Temporal quantity  $a$  with the activity set  $T_a \subseteq \mathcal{T}$  describes the changes of properties of nodes and links:

$$a = \begin{cases} a'(t) & t \in T_a \\ \text{undefined} & t \in \mathcal{T} \setminus T_a \end{cases}$$

Temporal quantities allow longitudinal approach instead of time slices.

$$\begin{aligned} a &= [(1, 5, 2), (6, 8, 1), (11, 12, 3), (14, 16, 2), (17, 18, 5), (19, 20, 1)] \\ b &= [(2, 3, 4), (4, 7, 3), (9, 10, 2), (13, 15, 5), (16, 21, 1)] \end{aligned}$$

The following are the sum  $s = a + b$  and the product  $p = a \cdot b$  of temporal quantities  $a$  and  $b$  over combinatorial semiring.

$$\begin{aligned} s &= [(1, 2, 2), (2, 3, 6), (3, 4, 2), (4, 5, 5), (5, 6, 3), (6, 7, 4), (7, 8, 1), \\ &\quad (9, 10, 2), (11, 12, 3), (13, 14, 5), (14, 15, 7), (15, 16, 2), (16, 17, 1), \\ &\quad (17, 18, 6), (18, 19, 1), (19, 20, 2), (20, 21, 1)] \\ p &= [(2, 3, 8), (4, 5, 6), (6, 7, 3), (14, 15, 10), (17, 18, 5), (19, 20, 1)] \end{aligned}$$

They are visually displayed at the bottom half of figures on the following slides.



# Addition and multiplication of temporal quantities

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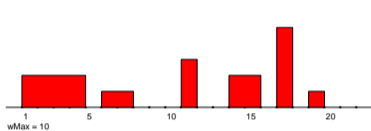
Algorithms

Results

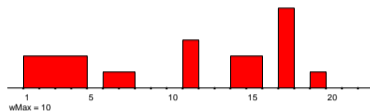
Conclusions

References

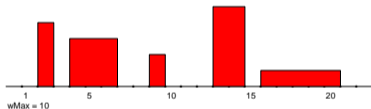
$a$  :



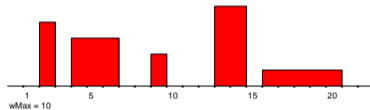
$a$  :



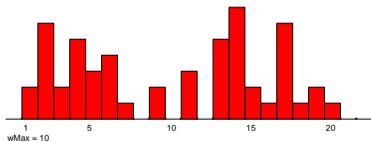
$b$  :



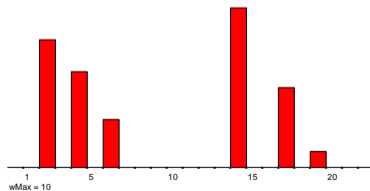
$b$  :



$a + b$  :



$a \cdot b$  :







# Core decomposition

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```
1 CoreDecomposition ( $\mathcal{N}$ ):
2  $C = V$ 
3  $k = 1$ 
4 while  $C \neq \emptyset$ :
5     while  $\exists u \in C \ni: \text{deg}(u) < k$ :
6         for  $v \in N(u, C)$ :
7              $\text{deg}(v) = \text{deg}(v) - 1$ 
8          $C = C \setminus v$ 
9          $\text{core}(u) = k - 1$ 
10     $k = k + 1$ 
```

We extended this algorithm to ordinary temporal cores and temporal  $p_S$ -cores.  
For programs in Python see [GitHub/Bavla/Graph](#).



# Algorithm for ordinary temporal cores

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```
1  TemporalCores( $\mathcal{N}$ ):
2   $D = \{u: [\text{triples}(\text{start}, \text{finish}, \text{deg})]\}$ 
3   $\text{CoreHierarchy} = \{u: [\text{triples with deg} = 0]\}$ 
4   $D = (D.\text{filter}(\text{deg} > 0)).\text{remove}(\text{empty triples})$ 
5   $D_{\min} = \{u: \text{min deg}\}$ 
6  while  $D$  not empty:
7       $(d_{\min}, u) = (\text{deg}, u) \ni: (u, \text{deg}) \in D_{\min} \wedge \text{deg is min deg}$ 
8       $\text{core} = [\text{triples from } D[u] \ni: \text{deg}[u] \text{ from triple is equal to } d_{\min}]$ 
9       $\text{CoreHierarchy}[u].\text{add}(\text{core})$ 
10      $\text{change} = \text{core}.\text{set}(\text{deg} = -1)$ 
11      $D[u] = D[u].\text{add}(\text{change}).\text{cutAt}(d_{\min}) \setminus \setminus \text{value} \geq d_{\min}$ 
12     for  $l$  in  $\mathcal{N}.\text{star}(u)$ :
13          $v = \text{other end-node of } l$ 
14         if not  $v$  in  $D$ : continue
15          $\text{changeLink} = l.\text{intersection}(\text{change}).\text{set}(\text{deg} = -1)$ 
16         if  $\text{changeLink}$  empty: continue
17          $\text{diff} = D[v].\text{add}(\text{changeLink}).\text{cutAt}(0) \setminus \setminus \text{value} \geq 0$ 
18          $D[v] = \text{diff}.\text{set}(\text{max}(\text{currentValue}, d_{\min}))$ 
19         if  $D[v]$  is empty:
20             delete  $D[v], D_{\min}[v]$ 
21         else:
22              $D_{\min}[v] = \text{triple} \in D[v] \text{ with min deg}$ 
23     if  $D[u]$  empty:
24         delete  $D[u], D_{\min}[u]$ 
25     else:
26          $D_{\min}[u] = \text{triple} \in D[u] \text{ with min deg}$ 
27 return  $\text{CoreHierarchy}$ 
```



# Algorithm for $p_S$ temporal cores

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```
1 PSTemporalCores( $\mathcal{N}$ ):
2  $D = \{u: [\text{triples}(\text{start}, \text{finish}, \text{weightSum})]\}$ 
3  $\text{Core} = \{u: [\text{triples with weightSum} = 0]\}$ 
4  $D = (D.\text{filter}(\text{weightSum} > 0)).\text{remove}(\text{empty triples})$ 
5  $D_{\min} = \{u: \text{min weightSum}\}$ 
6 while  $D$  not empty:
7    $(d_{\min}, u) = (\text{weightSum}, u) \ni D_{\min} \wedge \text{weightSum}$  is min weightSum
8    $\text{core} = [\text{triples from } D[u] \ni \text{weightSum}[u]$  from triple is equal to  $d_{\min}$ ]
9   if  $\text{core}$  not empty:
10     $\text{Core}[u].\text{add}(\text{core})$ 
11     $\text{change} = \text{core}.\text{set}(\text{weightSum} = -\text{weightSum})$ 
12     $D[u] = D[u].\text{add}(\text{change}).\text{cutAt}(d_{\min}) \setminus \setminus \text{value} \geq d_{\min}$ 
13    for  $l$  in  $\mathcal{N}.\text{star}(u)$ :
14       $v = \text{other end-node of } l$ 
15      if not  $v$  in  $D$ : continue
16       $\text{changeLink} = l.\text{intersection}(\text{change}).\text{set}(\text{weightSum} = -\text{weightSum})$ 
17      if  $\text{changeLink}$  empty: continue
18       $\text{diff} = D[v].\text{add}(\text{changeLink}).\text{cutAt}(0) \setminus \setminus \text{value} \geq 0$ 
19       $D[v] = \text{diff}.\text{set}(\text{max}(\text{currentValue}, d_{\min}))$ 
20      if  $D[v]$  is empty:
21        delete  $D[v], D_{\min}[v]$ 
22      else:
23         $D_{\min}[v] = \text{triple} \in D[v]$  with min weightSum
24    if  $D[u]$  is empty:
25      delete  $D[u], D_{\min}[u]$ 
26    else:
27       $D_{\min}[u] = \text{triple} \in D[u]$  with min weightSum
28 return  $\text{Core}$ 
```



# Artificial example

all weights  $w = 1$

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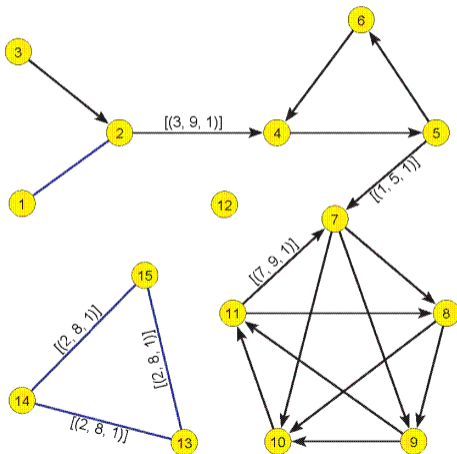
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# Artificial example

all weights  $w = 1$

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Node	Degree	Core number
<b>1</b>	(1, 9, 1)	(1, 9, 1)
<b>2</b>	(1, 3, 2), (3, 9, 3)	(1, 9, 1)
<b>3</b>	(1, 9, 1)	(3, 9, 1)
<b>4</b>	(1, 3, 2), (3, 9, 3)	(1, 9, 2)
<b>5</b>	(1, 5, 3), (5, 9, 2)	(1, 9, 2)
<b>6</b>	(1, 9, 2)	(1, 9, 2)
<b>7</b>	(1, 5, 4), (5, 7, 3), (7, 9, 4)	(1, 7, 3), (7, 9, 4)
<b>8</b>	(1, 9, 4)	(1, 7, 3), (7, 9, 4)
<b>9</b>	(1, 9, 4)	(1, 7, 3), (7, 9, 4)
<b>10</b>	(1, 9, 4)	(1, 7, 3), (7, 9, 4)
<b>11</b>	(1, 7, 3), (7, 9, 4)	(1, 7, 3), (7, 9, 4)
<b>12</b>	(1, 9, 0)	(1, 9, 0)
<b>13</b>	(1, 2, 0), (2, 8, 2), (8, 9, 0)	(1, 2, 0), (2, 8, 2), (8, 9, 0)
<b>14</b>	(1, 2, 0), (2, 8, 2), (8, 9, 0)	(1, 2, 0), (2, 8, 2), (8, 9, 0)
<b>15</b>	(1, 2, 0), (2, 8, 2), (8, 9, 0)	(1, 2, 0), (2, 8, 2), (8, 9, 0)



# Artificial example

## different weights

### Temporal cores

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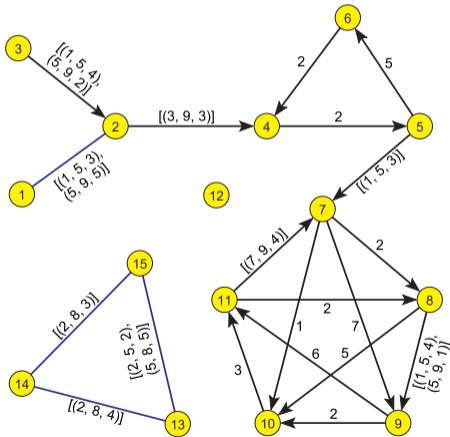
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# Artificial example

different weights

## Temporal cores

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Node	Degree	Core number
<b>1</b>	(1, 5, 3), (5, 9, 5)	(1, 5, 3), (5, 9, 5)
<b>2</b>	(1, 3, 7), (3, 9, 10)	(1, 5, 4), (5, 9, 5)
<b>3</b>	(1, 5, 4), (5, 9, 2)	(1, 5, 4), (5, 9, 2)
<b>4</b>	(1, 3, 4), (3, 9, 7)	(1, 5, 4), (5, 9, 5)
<b>5</b>	(1, 5, 10), (5, 9, 7)	(1, 9, 5)
<b>6</b>	(1, 9, 7)	(1, 9, 5)
<b>7</b>	(1, 5, 13), (5, 7, 10), (7, 9, 14)	(1, 9, 10)
<b>8</b>	(1, 5, 13), (5, 9, 10)	(1, 9, 10)
<b>9</b>	(1, 5, 19), (5, 9, 16)	(1, 9, 10)
<b>10</b>	(1, 9, 11)	(1, 9, 10)
<b>11</b>	(1, 7, 11), (7, 9, 15)	(1, 9, 10)
<b>12</b>	(1, 9, 0)	(1, 9, 0)
<b>13</b>	(1, 2, 0), (2, 5, 6), (5, 8, 9), (8, 9, 0)	(1, 2, 0), (2, 5, 5), (5, 8, 7), (8, 9, 0)
<b>14</b>	(1, 2, 0), (2, 8, 7), (8, 9, 0)	(1, 2, 0), (2, 5, 5), (5, 8, 7), (8, 9, 0)
<b>15</b>	(1, 2, 0), (2, 5, 5), (5, 8, 8), (8, 9, 0)	(1, 2, 0), (2, 5, 5), (5, 8, 7), (8, 9, 0)



# Reuters terror news network

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Obtained from the CRA (Centering Resonance Analysis) networks produced by Steve Corman and Kevin Dooley at Arizona State University.

Based on all the stories released during 66 consecutive days by the news agency Reuters concerning the September 11 attack on the U.S., beginning at 9:00 AM EST 9/11/01.

Nodes: important words (terms),  $n = 13332$

Links: two nodes appear in the same utterance,  $m = 243447$ , undirected, weight is equal to the frequency of appearance, 50859 of them have the weight larger than 1. No loops.

Data available at: [Terror](#).

Example: induced subnetwork on 50 most active nodes.





# Reuters terror news network

## Temporal degrees

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Node	Degree
1	(1, 2, 5), (2, 3, 6), (3, 4, 3), (4, 5, 5), (5, 6, 4), (6, 8, 3), (8, 10, 5), (10, 11, 3), (11, 13, 2), (13, 16, 3), (16, 17, 4), (17, 18, 5), (18, 19, 3), (19, 21, 1), (21, 22, 2), (22, 23, 1), (23, 24, 4), (24, 25, 1), (25, 29, 3), (29, 31, 2), (31, 33, 3), (33, 34, 1), (34, 36, 3), (36, 37, 2), (37, 39, 3), (39, 40, 4), (40, 41, 2), (41, 42, 0), (42, 43, 3), (43, 44, 2), (44, 45, 3), (45, 46, 1), (46, 47, 2), (47, 48, 3), (48, 49, 0), (49, 50, 4), (50, 51, 1), (51, 52, 2), (52, 53, 1), (53, 54, 0), (54, 58, 2), (58, 59, 3), (59, 60, 2), (60, 61, 4), (61, 62, 0), (62, 64, 2), (64, 65, 1), (65, 67, 2)
2	(1, 2, 27), (2, 3, 29), ..., (63, 64, 2), (64, 65, 0), (66, 67, 0)
...	
50	(1, 2, 3), (2, 3, 2), (3, 5, 1), (5, 8, 0), (8, 10, 1), (10, 11, 2), (11, 12, 1), (12, 15, 0), (15, 16, 3), (16, 17, 1), (17, 19, 0), (19, 20, 1), (20, 21, 2), (21, 22, 0), (22, 24, 1), (24, 26, 0), (26, 27, 2), (27, 28, 0), (28, 29, 1), (29, 31, 0), (31, 32, 1), (32, 33, 0), (33, 35, 1), (35, 37, 0), (37, 38, 1), (38, 42, 0), (43, 44, 2), (44, 49, 0), (49, 50, 2), (51, 57, 0), (58, 61, 0), (61, 62, 1), (62, 67, 0)



# Reuters terror news network

## Temporal cores

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Node	Core number
1	(1, 2, 4), (2, 3, 5), (3, 5, 3), (5, 6, 4), (6, 8, 3), (8, 10, 4), (10, 11, 3), (11, 14, 2), (14, 18, 3), (18, 19, 2), (19, 21, 1), (21, 22, 2), (22, 23, 1), (23, 24, 3), (24, 25, 1), (25, 28, 2), (28, 29, 3), (29, 33, 2), (33, 34, 1), (34, 38, 2), (38, 39, 3), (39, 41, 2), (41, 42, 0), (42, 45, 2), (45, 46, 1), (46, 47, 2), (47, 48, 3), (48, 49, 0), (49, 50, 3), (50, 51, 1), (51, 52, 2), (52, 53, 1), (53, 54, 0), (54, 57, 2), (57, 58, 1), (58, 59, 2), (59, 60, 1), (60, 61, 2), (61, 62, 0), (62, 64, 2), (64, 65, 1), (65, 67, 2)
2	(1, 3, 5), (3, 6, 4), (6, 7, 5), ..., (63, 64, 1), (64, 65, 0), (66, 67, 0)
...	
50	(1, 3, 2), (3, 5, 1), (5, 8, 0), (8, 10, 1), (10, 11, 2), (11, 12, 1), (12, 15, 0), (15, 16, 3), (16, 17, 1), (17, 19, 0), (19, 20, 1), (20, 21, 2), (21, 22, 0), (22, 24, 1), (24, 26, 0), (26, 27, 1), (27, 28, 0), (28, 29, 1), (29, 31, 0), (31, 32, 1), (32, 33, 0), (33, 35, 1), (35, 37, 0), (37, 38, 1), (38, 42, 0), (43, 44, 1), (44, 49, 0), (49, 50, 2), (51, 57, 0), (58, 61, 0), (61, 62, 1), (62, 67, 0)



# Reuters terror news network

Temporal cores of order at least 3 appear in the first 11 days and on 30th day

## Temporal cores

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	<b>Node</b>	<b>Core number (<math>\geq 3</math>)</b>
25	world	(1, 3, 5), (3, 10, 4)
2	attack	(1, 3, 5), (3, 6, 4), (6, 7, 5), (7, 10, 4), (11, 12, 4), (30, 31, 4)
9	washington	(1, 3, 5), (3, 6, 4), (6, 7, 5), (7, 10, 4), (11, 12, 4)
14	world_trade_ctr	(1, 3, 5), (3, 6, 4), (6, 7, 5), (30, 31, 4)
4	people	(1, 3, 5), (3, 6, 4), (6, 7, 5), (7, 8, 4)
21	pentagon	(1, 3, 5), (3, 4, 4), (5, 6, 4), (6, 7, 5)
7	new_york	(1, 3, 5), (3, 6, 4), (6, 7, 5), (30, 31, 4)
8	pres_bush	(1, 3, 5), (3, 6, 4), (6, 7, 5), (7, 10, 4), (11, 12, 4)
10	official	(1, 3, 5), (3, 4, 4), (5, 6, 4), (6, 7, 5)
43	tower	(1, 3, 5), (3, 4, 4), (6, 7, 5)
34	time	(1, 3, 5), (3, 4, 4), (5, 6, 4), (7, 8, 4)
18	city	(1, 3, 5), (3, 4, 4)
20	tuesday	(1, 3, 5), (3, 7, 4)
13	plane	(1, 3, 5), (3, 7, 4)
15	security	(1, 2, 4), (2, 3, 5), (5, 6, 4)
1	united_states	(1, 2, 4), (2, 3, 5), (5, 6, 4), (8, 10, 4)
19	war	(1, 2, 4), (2, 3, 5), (5, 8, 4)
29	worker	(1, 2, 4), (2, 3, 5)
47	wednesday	(2, 3, 5), (3, 4, 4), (8, 10, 4)
12	military	(1, 2, 4), (5, 6, 4), (30, 31, 4)
5	afghanistan	(1, 3, 4), (5, 6, 4), (6, 7, 5), (8, 10, 4), (30, 31, 4)



# Reuters terror news network

Temporal cores of order at least 3 appear in the first 11 days and on 30th day

## Temporal cores

V. Batagelj,  
M. Cerinšek

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References

	<b>Node</b>	<b>Core number (<math>\geq 3</math>)</b>
6	bin_laden	(1, 4, 4), (5, 6, 4), (6, 7, 5), (7, 10, 4), (11, 12, 4)
36	strike	(2, 3, 4), (5, 6, 4), (6, 7, 5), (30, 31, 4)
28	week	(5, 6, 4), (6, 7, 5), (8, 10, 4), (11, 12, 4)
48	nation	(1, 3, 4), (5, 6, 4)
40	terrorist	(1, 3, 4), (6, 7, 4)
17	country	(1, 3, 4), (5, 10, 4)
23	government	(1, 3, 4), (5, 6, 4)
30	office	(1, 3, 4)
24	leader	(1, 4, 4), (6, 10, 4)
49	police	(2, 4, 4), (5, 6, 4)
31	group	(2, 3, 4), (6, 7, 4)
42	pakistan	(2, 3, 4), (5, 7, 4)
32	air	(2, 3, 4), (5, 6, 4)
27	day	(2, 3, 4), (5, 6, 4)
35	hijack	(2, 3, 4)
26	terrorism	(2, 3, 4)
38	flight	(2, 3, 4)
39	tell	(2, 3, 4)
16	american	(2, 3, 4)
41	airport	(2, 3, 4)
45	new	(2, 3, 4)
22	force	(5, 6, 4)



# Reuters terror news network

## Temporal $p_5$ -cores

### Temporal cores

V. Batagelj,  
M. Cerinšek

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References

	Node	$p_5$ -core number ( $\geq 20$ )
2	attack	(1, 3, 86), (3, 4, 44), (4, 5, 36), (5, 6, 66), (6, 7, 47), (7, 8, 22), (8, 9, 21), (9, 10, 24), (10, 11, 22), (11, 12, 24), (15, 16, 23), (18, 19, 20), (27, 28, 23)
7	new_york	(1, 2, 101), (2, 3, 86), (3, 4, 42), (4, 5, 35), (5, 6, 66), (6, 7, 47), (8, 9, 21), (9, 10, 24), (10, 11, 22), (11, 12, 24), (15, 16, 23), (18, 19, 20)
14	world_trade_c	(1, 2, 101), (2, 3, 86), (3, 4, 44), (4, 5, 35), (5, 6, 66), (6, 7, 47), (8, 9, 20), (9, 10, 21), (10, 11, 22), (11, 12, 24), (15, 16, 23), (18, 19, 20)
9	washington	(1, 2, 80), (2, 3, 61), (3, 4, 27), (4, 5, 28), (5, 6, 66), (6, 7, 47), (8, 9, 21), (9, 10, 24), (10, 11, 22), (11, 12, 24), (15, 16, 23), (18, 19, 20)
21	pentagon	(1, 3, 86), (3, 4, 44), (4, 5, 32), (5, 6, 66), (6, 7, 47), (8, 9, 20), (9, 10, 21), (10, 11, 22), (11, 12, 24), (15, 16, 23), (18, 19, 20)
1	united_states	(1, 2, 86), (2, 3, 71), (3, 4, 34), (4, 5, 29), (5, 6, 50), (6, 7, 47), (7, 8, 22), (15, 16, 23), (18, 19, 23), (27, 28, 23)
28	week	(5, 6, 35), (6, 7, 27), (7, 8, 22), (8, 9, 21), (9, 10, 24), (10, 11, 22), (11, 12, 24)
4	people	(1, 2, 48), (2, 3, 52), (3, 4, 28), (4, 5, 32), (5, 6, 29), (6, 7, 34), (18, 19, 20)
12	military	(1, 2, 25), (2, 3, 42), (5, 6, 26), (15, 16, 23), (18, 19, 23), (27, 28, 23)
5	afghanistan	(1, 2, 22), (2, 3, 28), (5, 6, 29), (6, 7, 21), (15, 16, 23), (27, 28, 23)
6	bin_laden	(1, 2, 22), (2, 3, 28), (3, 4, 20), (5, 6, 29), (6, 7, 21), (18, 19, 20)
10	official	(1, 2, 40), (2, 3, 54), (3, 4, 34), (5, 6, 29), (6, 7, 36), (18, 19, 23)
43	tower	(1, 2, 101), (2, 3, 72), (3, 4, 41), (4, 5, 32), (5, 6, 38), (6, 7, 32)
35	hijack	(1, 2, 67), (2, 3, 86), (3, 4, 44), (4, 5, 28), (5, 6, 50), (6, 7, 34)
13	plane	(1, 3, 86), (3, 4, 44), (4, 5, 32), (5, 6, 50), (6, 7, 34)



# Reuters terror news network

## Temporal $p_5$ -cores

### Temporal cores

V. Batagelj,  
M. Cerinšek

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	Node	$p_5$ -core number ( $\geq 20$ )
20	tuesday	(1, 3, 86), (3, 4, 44), (4, 5, 36), (5, 6, 66), (6, 7, 47)
3	taliban	(2, 3, 28), (6, 7, 20), (15, 16, 23), (27, 28, 23)
36	strike	(2, 3, 29), (5, 6, 29), (18, 19, 22), (27, 28, 23)
17	country	(1, 2, 24), (2, 3, 31), (5, 6, 26), (18, 19, 20)
8	pres_bush	(1, 2, 48), (2, 3, 44), (5, 6, 29), (6, 7, 21)
41	airport	(1, 2, 25), (2, 3, 44), (4, 5, 25), (5, 6, 24)
15	security	(1, 2, 25), (2, 3, 30), (5, 6, 24)
16	american	(1, 2, 48), (2, 3, 30), (5, 7, 20)
18	city	(1, 2, 60), (2, 3, 52), (3, 4, 22)
25	world	(1, 2, 34), (2, 3, 44), (18, 19, 20)
27	day	(1, 2, 21), (2, 3, 36), (5, 6, 20)
32	air	(2, 3, 34), (5, 6, 29), (27, 28, 23)
38	flight	(1, 2, 25), (2, 3, 52), (4, 5, 20)
48	nation	(1, 2, 31), (2, 3, 38), (5, 6, 23)
40	terrorist	(1, 2, 40), (2, 3, 29)
19	war	(2, 3, 34), (5, 6, 29)
23	government	(1, 2, 28), (2, 3, 36)
46	buildng	(1, 2, 34), (2, 3, 44)
30	office	(1, 2, 34), (2, 3, 20)



# Reuters terror news network

## Temporal $p_5$ -cores

### Temporal cores

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	<b>Node</b>	<b><math>p_5</math>-core number (<math>\geq 20</math>)</b>
26	terrorism	(5, 6, 20)
29	worker	(1, 2, 24)
31	group	(2, 3, 26)
34	time	(2, 3, 36)
22	force	(5, 6, 26)
24	leader	(1, 2, 22)
42	pakistan	(5, 6, 29)
44	bomb	(1, 2, 23)
45	new	(2, 3, 30)
47	wednesday	(2, 3, 52)
49	police	(2, 3, 20)



# Max $p_5$ -core numbers by days from the event

## Temporal cores

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M. Cerinšek

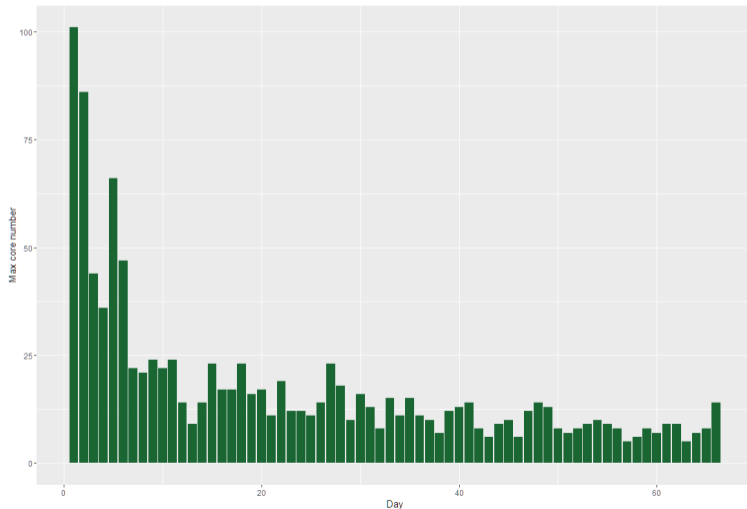
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# Franzosi's Violence network

## Temporal cores

V. Batagelj,  
M. Cerinšek

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Roberto Franzosi collected from the journal news in the period January 1919 - December 1922 information about the different types of interactions between political parties and other groups of people in Italy. The violence network contains only the data about violent actions and counts the number of interactions per month.

Nodes: groups of people,  $n = 29$

Links: violent interactions,  $m = 105$

For details see:

Franzosi, R., 1997. Mobilization and CounterMobilization Processes: From the Red Years (1919-20) to the Black Years (1921-22) in Italy.

Franzosi, R., 1997. A New Methodological Approach to the Study of Narrative Data. Theory and Society, 26(2-3), 275-304



# Violence network

Core number  $\geq 3$

## Temporal cores

V. Batagelj,  
M. Cerinšek

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References

<b>Node</b>	<b>Core number (<math>\geq 3</math>)</b>
16 workers	(29, 30, 3), (33, 34, 3), (39, 41, 3)
1 undefined	(29, 30, 3), (39, 40, 3)
2 ?	(31, 32, 3), (33, 34, 3), (40, 41, 3)
3 people	(31, 32, 3), (33, 34, 3), (39, 40, 3)
4 police	(31, 32, 3), (33, 34, 3), (40, 41, 3)
21 catholics	(33, 34, 3)
7 fascists	(29, 30, 3), (31, 32, 3), (33, 34, 3), (39, 41, 3)
8 communists	(29, 30, 3)
10 socialists	(31, 32, 3), (40, 41, 3)



# Violence network

Core number  $\geq 2$

## Temporal cores

V. Batagelj,  
M. Cerinšek

Definitions

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References

Node	Core number ( $\geq 2$ )
1 undefined	(15, 16, 2), (17, 18, 2), (25, 29, 2), (29, 30, 3), (31, 32, 2), (38, 39, 2), (39, 40, 3), (41, 44, 2), (45, 46, 2), (48, 49, 2)
2 ?	(14, 16, 2), (17, 18, 2), (28, 29, 2), (31, 32, 3), (32, 33, 2), (33, 34, 3), (34, 35, 2), (40, 41, 3)
3 people	(16, 18, 2), (23, 24, 2), (25, 26, 2), (28, 30, 2), (31, 32, 3), (33, 34, 3), (35, 37, 2), (39, 40, 3), (41, 43, 2), (48, 49, 2)
4 police	(11, 12, 2), (14, 20, 2), (21, 23, 2), (29, 31, 2), (31, 32, 3), (32, 33, 2), (33, 34, 3), (34, 37, 2), (38, 40, 2), (40, 41, 3)
5 land owners	(15, 16, 2), (17, 20, 2), (29, 30, 2), (36, 37, 2), (38, 40, 2), (42, 43, 2)
7 fascists	(11, 12, 2), (16, 17, 2), (19, 20, 2), (21, 24, 2), (25, 29, 2), (29, 30, 3), (30, 31, 2), (31, 32, 3), (32, 33, 2), (33, 34, 3), (34, 37, 2), (38, 39, 2), (39, 41, 3), (41, 44, 2), (45, 46, 2), (48, 49, 2)
8 communists	(28, 29, 2), (29, 30, 3), (31, 33, 2), (35, 37, 2), (43, 44, 2)
9 workers (agr)	(15, 16, 2), (17, 20, 2), (28, 30, 2), (31, 32, 2), (33, 35, 2), (38, 43, 2), (45, 46, 2)
10 socialists	(11, 12, 2), (16, 18, 2), (19, 20, 2), (22, 23, 2), (25, 26, 2), (27, 30, 2), (31, 32, 3), (33, 37, 2), (38, 40, 2), (40, 41, 3), (41, 42, 2)
12 war affected	(35, 36, 2), (39, 40, 2)
13 protesters	(15, 16, 2), (21, 22, 2), (29, 30, 2), (31, 32, 2), (38, 40, 2)



# Violence network

Core number  $\geq 2$

## Temporal cores

V. Batagelj,  
M. Cerinšek

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References

Node	Core number ( $\geq 2$ )
16 workers	(11, 12, 2), (14, 18, 2), (19, 20, 2), (21, 24, 2), (25, 26, 2), (27, 29, 2), (29, 30, 3), (30, 33, 2), (33, 34, 3), (34, 37, 2), (38, 39, 2), (39, 41, 3), (41, 44, 2), (45, 46, 2)
17 the right	(17, 18, 2), (41, 42, 2)
19 populars	(41, 42, 2)
20 students	(17, 18, 2)
21 catholics	(33, 34, 3)
25 republicans	(26, 27, 2)
26 thugs	(29, 30, 2)
27 prisoners/arrested	(40, 41, 2)



# Violence network

Temporal  
cores

V. Batagelj,  
M. Cerinšek

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Node		$p_S$ -ore number ( $\geq 10$ )
16	workers	(1, 2, 27), (10, 11, 11), (14, 15, 27), (16, 17, 11), (17, 18, 17), (18, 19, 12), (22, 23, 17), (25, 26, 11), (27, 28, 18), (28, 29, 16), (29, 30, 53), (30, 31, 56), (31, 32, 51), (32, 33, 30), (33, 34, 17), (34, 35, 71), (35, 36, 76), (36, 37, 53), (37, 38, 11), (38, 39, 23), (39, 40, 54), (40, 41, 13), (41, 42, 174), (42, 43, 25), (43, 44, 20), (45, 46, 15), (46, 47, 25)
10	socialists	(10, 11, 10), (12, 13, 29), (27, 28, 30), (28, 29, 31), (29, 30, 64), (30, 31, 29), (31, 32, 17), (32, 33, 14), (33, 34, 24), (34, 35, 38), (35, 36, 23), (36, 37, 26), (37, 38, 13), (38, 39, 19), (39, 40, 54), (45, 46, 13)
4	police	(1, 2, 36), (6, 7, 15), (10, 11, 24), (12, 13, 29), (14, 15, 27), (15, 16, 13), (16, 17, 24), (17, 18, 17), (18, 19, 12), (22, 23, 17), (31, 32, 17)
7	fascists	(25, 26, 11), (27, 28, 30), (28, 29, 31), (29, 30, 64), (30, 31, 56), (31, 32, 51), (32, 33, 30), (33, 34, 24), (34, 35, 71), (35, 36, 76), (36, 37, 53), (37, 38, 13), (38, 39, 23), (39, 40, 54), (40, 41, 13), (41, 42, 174), (42, 43, 25), (43, 44, 20), (45, 46, 15), (46, 47, 25)
9	workers (agr)	(10, 11, 24), (16, 17, 24), (28, 29, 16), (30, 31, 13), (36, 37, 11), (39, 40, 15), (43, 44, 10)
1	undefined	(25, 26, 11), (27, 28, 12), (28, 29, 16), (41, 42, 133), (45, 46, 11)
8	communists	(29, 30, 13), (30, 31, 10), (31, 32, 12)
13	protesters	(6, 7, 15), (15, 16, 13), (16, 17, 20)
12	war affected	(1, 2, 36)
3	people	(28, 29, 12)



# SN5 network

## Temporal cores

V. Batagelj,  
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Data from Web of Science ("social network\*" AND S0=(Social networks)) plus most frequently cited works plus around 100 SNA researchers. Collected in December 2007 for the 2008 Vizards session.

We analyze the works  $\times$  authors network **WA** restricted to works with a complete description ( $DC > 0$ ):  $|\mathbf{W}| = 7950$ ,  $|\mathbf{A}| = 12458$  and  $m = 19488$ .

Using the publication years the network **WA** was expanded to a temporal network with cumulative weights.

Normalization "by rows":  $\mathbf{N} = n(\mathbf{WA})$ .

Normalized coauthorship network:  $\mathbf{Ct} = \mathbf{N}^T * \mathbf{N}$ .



# SN5 Coauthorship

Temporal  
cores

V. Batagelj,  
M. Cerinšek

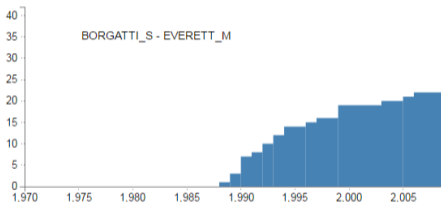
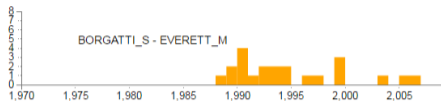
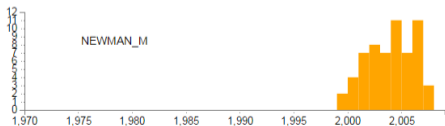
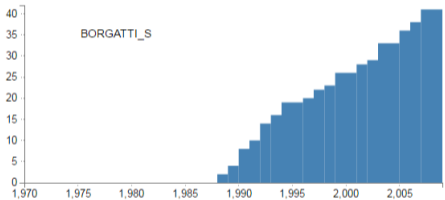
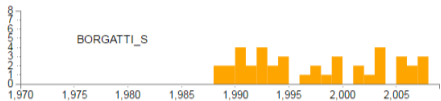
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# SN5 cumulative coauthorship $p_S$ cores $\geq 3$

## Temporal cores

V. Batagelj,  
M. Cerinšek

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```
20 : BORGATTI_S [(1991, 1992, 3.1667), (1992, 1993, 4.1667), (1993, 1994, 5.1667),
(1994, 1996, 6.1667), (1996, 1997, 6.6667), (1997, 1999, 7.1667), (1999, 2003, 8.6667),
(2003, 2005, 8.7917), (2005, 2006, 9.2917), (2006, 2009, 9.7917)]
3169 : EVERETT_M [(1991, 1992, 3.1667), (1992, 1993, 4.1667), (1993, 1994, 5.1667),
(1994, 1996, 6.1667), (1996, 1997, 6.6667), (1999, 2003, 8.6667), (2003, 2005, 8.7917),
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317 : BERNARD_H [(1990, 1991, 3.0244), (1991, 1995, 3.1494), (1995, 1997, 3.3094),
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4860 : METZKE_C [(2003, 2005, 3.0), (2005, 2006, 3.2222), (2006, 2009, 3.6667)]
3125 : SHELLEY_G [(2006, 2009, 3.4767)]
1673 : MCCARTY_C [(2006, 2009, 3.4767)]
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79 : NEWMAN_M [(2005, 2009, 3.0)]
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```





# Conclusions

## Temporal cores

V. Batagelj,  
M. Cerinšek

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References

- 1 Improve the complexity of the algorithm
- 2 Extend the algorithm to generalized temporal cores
- 3 Find user friendly presentations of results
- 4 Compare with the streaming core algorithms

Temporal Quantities – a Python 3 library for temporal network analysis [TQ](#) / [Graph](#).

### **Acknowledgements.**

The first author was partially sponsored by Slovenian Research Agency (ARRS) – projects J7-8279 and J1-6720, and grant P1-0294.

The second author was partially sponsored by Slovenian Research Agency (ARRS) – projects Z7-7614 (B) and L7-5554.



# References I

## Temporal cores

V. Batagelj,  
M. Cerinšek








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References

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## References II

### Temporal cores

V. Batagelj,  
M. Cerinšek

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